Exploiting problem structure in optimization with GeNIOS.jl (GEneralized Newtown Inexact Operator Splitting solver)

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Huber fitting (as a quadratic program)

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Markowitz Portfolio Optimization

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Harder to take advantage of problem data structure in this form

Example: robust (Huber) regression Quadratic program formulation has ~5x the variables!

minimize
$$\sum_{i=1}^{N} \ell^{\text{hub}}(a_i^T x - b_i) + \lambda_1 ||x||_1,$$

$$\ell^{\text{hub}}(w) = \begin{cases} w^2 & |w| \le 1\\ 2|w| - 1 & |w| > 1. \end{cases}$$

Huber robust regression

minimize $r^T r + 2\mathbf{1}^T (s+t) + \lambda_1 q$ subject to Ax - r - s + t = b $-q \le x \le q$ $0 \le s, t$

Equivalent quadratic program

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minimize
$$\sum_{i=1}^{N} \ell^{\text{hub}}(a_i^T x - b_i) + \lambda_1 \|x\|_1$$

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 $f(x), \nabla f(x), v \mapsto \hat{\nabla}^2 f(x)v$ $g(z), \text{ prox}_{g/\rho}(z)$ minimize f(x) + g(z)subject to Mx - z = c,



minimize $(1/2)x^T P x + q^T x$ subject to $l \leq Ax \leq u$,

Quadratic Programs

 $f(x), \nabla f(x), v \mapsto \hat{\nabla}^2 f(x)v$ $g(z), \hat{\mathbf{prox}}_{g/o}(z)$ minimize f(x) + g(z)subject to Mx - z = c,



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Quadratic Programs

minimize subject to A



$$(1/2)x^T P x + q^T x$$
$$Ax - z = c$$
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Conic Programs



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 - Approximate prox, $\operatorname{aprox}_{\varrho}(v)$, operator for 2nd subproblem
- Then we **inexactly solve** these approximate subproblems



Back to robust (Huber) regression Quadratic program formulation has ~5x the variables!

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$$\sum_{i=1}^{N} \ell^{\text{hub}}(a_i^T x - b_i) + \lambda_1 ||x||_1,$$

$$\ell^{\text{hub}}(w) = \begin{cases} w^2 & |w| \le 1\\ 2|w| - 1 & |w| > 1. \end{cases}$$

Huber robust regression

Can try solving QP with/without inexact subproblem solves

minimize $r^T r + 2\mathbf{1}^T (s+t) + \lambda_1 q$ subject to Ax - r - s + t = b $-q \leq x \leq q$ $0 \leq s, t$

Equivalent quadratic program

MLSolver only needs loss & regularization

Huber problem: min ∑ fʰuʰ(ai™x – bi) + λ||x||ı $f(x) = abs(x) <= 1 ? 0.5*x^2 : abs(x) - 0.5$ $\lambda 1 = \lambda$ $\lambda 2 = 0.0$ solver = GeNIOS_MLSolver(f, $\lambda 1$, $\lambda 2$, A, b)

minimize $\sum \ell(a_i^T x - b_i) + \lambda_1 ||x||_1 + (1/2)\lambda_2 ||x||_2^2$,

res = solve!(solver; options=GeNIOS_SolverOptions(use_dual_gap=false))

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Supply conjugate -> can use duality gap convergence criterion

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QPSolver just needs problem data mats/vecs (Can be supplied through JuMP)

- minimize $(1/2)x^T P x + q^T x$ subject to $l \leq Ax \leq u$,

solver = GeNIOS_QPSolver(P, q, A, l, u) res = solve!(solver)



QPSolver just needs problem data mats/vecs (Can be supplied through JuMP)

- subject to $l \leq Ax \leq u$,

minimize $(1/2)x^T P x + q^T x$





MLSolver gives 35x-60x speedup





Sparsification slows mat-vec-products (explains wider gap bt linsys & full solve)

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Sounds like a lot of work...



Markowitz Portfolio Optimization i.e., maximize risk-adjusted return of a long-only portfolio

- subject to $\mathbf{1}^T x = 1$

minimize $-\mu^T x + (\gamma/2) x^T \Sigma x$ $x \ge 0,$

Markowitz Portfolio Optimization i.e., maximize risk-adjusted return of a long-only portfolio

- x > 0,
- minimize $-\mu^T x + (\gamma/2) x^T \Sigma x$ subject to $\mathbf{1}^T x = 1$ • Covariance is <u>diagonal + low rank</u>: $\Sigma = D + FF^T$ where $F \in \mathbb{R}^{n \times k}$, $k \ll n$

Markowitz Portfolio Optimization i.e., maximize risk-adjusted return of a long-only portfolio

- minimize $-\mu^T x + (\gamma/2) x^T \Sigma x$ subject to $\mathbf{1}^T x = 1$ x > 0,
- Covariance is <u>diagonal + low rank</u>: $\Sigma = D + FF^T$ where $F \in \mathbb{R}^{n \times k}$, $k \ll n$
- This is a quadratic program:
 - minimize $(1/2)x^TPx + q^Tx$ subject to l < Ax < u,

How can we take advantage of structure? Traditional method: solve equivalent QP with n+k variables

- An equivalent QP w new variable $y \in \mathbf{R}^k$:
 - minimize $\gamma x^T D x + \gamma y^T y \mu^T x$ subject to $y = F^T x$ ${\bf 1}^T x = 1$
- $x \ge 0.$

- Pro: much faster solve

Con: burden on user to do reformulation (can be very complex in other cases!)



Julia way: multiple dispatch We can create fast mat-vec-products for constraint & obj matrices

using LinearMaps ## P = $\gamma * (F * F' + Diagonal(d))$ F_lm = LinearMap(F) P = \vec{P*(F_lm*F_lm' + Diagonal(d)) ## M = vcat(I, ones(1, n))M = vcat(LinearMap(I, n), ones(1, n))solver = GeNIOS_QPSolver(P, q, M, l, u; check_dims=false);

res = solve!(solver; options=GeNIOS_SolverOptions(eps_abs=1e-6));

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• But the solver allows us to solve an even more interesting formulation...

res = solve!(solver; options=GeNIOS_SolverOptions(eps_abs=1e-6));

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- Equivalent convex (non-QP) problem:
 - minimize $-\mu^T x$ subject to x - z
- where I_S is an indication function of the set $S = \{z \mid \mathbf{1}^T z = 1 \text{ and } z \ge 0\}$

$$x + (\gamma/2)x^T\Sigma x + I_S(z)$$

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- Z-subproblem is a projection onto S
 - can be solved by 1d root finding (read: really really fast)

$$\begin{aligned} x + (\gamma/2) x^T \Sigma x + I_S(z) \\ = 0, \end{aligned}$$

Generic problem can be specified directly

solver = GeNIOS_GenericSolver(f, grad_f!, Hf, # f(x)g, prox_g!, I, zeros(n) res = solve!(solver)

minimize f(x) + g(z)subject to Mx - z = c,

g(z) # M, C: MX - Z = C

Solve the original QP

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• How do they compare?

In short: the more structure used, the better And Julia lets us avoid error-prone reformulations



Thank you! Questions?

Code & documentation:

<u>https://github.com/tjdiamandis/GeNIOS.jl</u>

Relevant theory paper:

Frangella, Z., Zhao, S., Diamandis, T., Stellato, B., & Udell, M. (2023). On the (linear) convergence of Generalized Newton Inexact ADMM. arXiv preprint arXiv:2302.03863.

- Plan to expose more of the interface in JuMP, extend to more sets
- Email: tdiamand@mit.edu (or open an issue on GitHub)



Appendix

ADMM is a popular first-order method for constrained optimization

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- Core idea: split original problem (difficult) into 2+ easy problems
 - Repeatedly solve these problems & push solns together

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 - Repeatedly solve these problems & push solns together

• Very effective for large-scale, data-driven opt problems

Distributed optimization and statistical learning via the **alternating direction** method of multipliers

S Boyd, N Parikh, E Chu, B Peleato ... - ... and Trends® in ..., 2011 - nowpublishers.com ... review, we argue that the alternating direction method of multipliers is well suited to distributed convex optimization, and in particular to large-scale problems arising in statistics, ... ☆ Save 50 Cite Cited by 20291 Related articles All 43 versions ≫

$$x^{k+1} = \underset{x}{\operatorname{argmin}} \left(f(x) + \underset{z}{\operatorname{argmin}} \left(g(z) + \underset{z}{\operatorname{argmin}} \right) \right)$$
$$u^{k+1} = u^k + M x^{k+1} - u^k +$$

 $egin{aligned} &-(
ho/2)\|Mx-z^k-c+u^k\|_2^2\ &-(
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- But solving the x-subproblem can still be difficult / slow
 - Generally have to run an algorithm like L-BFGS at each iteration
- A (perhaps silly?) idea: replace f(x) with an easy approximation

 $egin{aligned} & (
ho/2) \| Mx - z^k - c + u^k \|_2^2 \ & (
ho/2) \| Mx^{k+1} - z - c + u^k \|_2^2 \ & z^{k+1} - c. \end{aligned}$

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- Idea 1: replace f(x) with the second order Taylor expansion around x^k
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 - Makes interface very easy (& can leverage auto diff)

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- Idea 2: solve this linear system inexactly
 - We solve with CG method, decreasing the tolerance at each iteration

What about the z-subproblem?

- Often can put the hard/time-consuming parts of the problem in f
- But theory tells us we can solve z-subproblem inexactly too!

- We won't discuss here
- But interesting future directions...